Searching

Blind search – keep moving randomly till reach goal, independent of the domain depth first, breadth-first, iterative-deepening

Heuristic – utilities domain specific features, best first A* i.e. Finding a word in a dictionary we know it starts with say Z so look there first

Local – if only finding goal state matters i.e. Situations, hill-climbing, simulate annealing, genetic algorithms

Blind Search

Tree Search – just go down the tree till hit sumthing good

Strategy's – Order of expansion

completeness – always find a solution if its there

time complexity – number of nodes generated / expanded how long will it take

space complexity – maximum number of nodes in memory at any one time optimality – does it always find the least time cost solution ie fastest

time and space complexity measured by

b = maximum branching factor

d = depth of best solution

m = maximum tree depth

Breadth first

work along each level of the tree expand shallowest unexpanded node fringe as FIFO queue

has to store all nodes in memory

Depth first

work down to the bottom of each branch expand deepest unexpanded nodes fringe as LIFO queue

Not optimal can get stuck on infinite paths Less memory than Breadth first

Iterative Deepening

effectively breadth first but take advantage of depth search by doing a depth search to iteratively increasing depths

Open and Closed nodes

Remember visited nodes to stop infinite loops

Heuristic Search

Use other known knowledge to search for solutions

Best first search – pick the seemingly best node among the search nodes based on an

evaluation function that works for the data.

- i.e. Search inwards towards a goal
- i.e. With a compass

A* search

cost (s->n) called g(n) cost (n->g) called h(n)

g(n) is path cost to n

h(n) is estimate of least cost path from n to g

f(n) = g(n) + h(n) estimated cost of cheapest solutions among search nodes select node where f(n) is lowest

depending on h(n) this can be complete and optimal if h(n) = < true cost(n->g)

for A^* if h(n) is admissible then its optimal

complete, unless infinite many nodes f=<f(g) time, exponential space, needs to store all nodes optimal, yes

Iterative Improvement

i.e. The chess game

Hill Climbing

trying to find the shallowest part of the sea using a probe step along if its increasing keep going, if it starts going down go back and try and find highest

can miss highest and home in on a smaller point

Simulated Annealing

escape from local maximums by allowing some bad moves, have a probability that you can move

used a lot as it allows solutions to hard problems to be effectively guessed / worked towards

$$P(x) \propto exp(-E(x) / T)$$

T slowly reaches so that it finds goal

Genetic algorithms

Pick 2 parents apply mutation bit flip with probability replace population with offspring

1 point crossover just cut and join the 2 strings somewhere

selection

give individuals chances and place these on a roulette wheel then spin the wheel n times to get n individuals

works on local search in uncertain conditions

MiniMax search

Max is player A, aims to pick best moves Min is player B, aims to pick best moves to minimise A's advantage

The faster the computer the deeper it can look

Cant look forever, if opponent looks 1 more than us then we have no chance, especially if the opponent notices this.

Optimise

take into account board size

Alpha-Beta pruning

Alpha lower bound on node evaluations (worst we can do)
Associate with max nodes
Never decreases
Beta represents the upper bound
Associate with min nodes
Never increases